

Last time: (Thm 3.28):

Let G be a graph and let $X, Y \subseteq V(G)$ be non-empty.

The minimum size of a set that separates X from Y

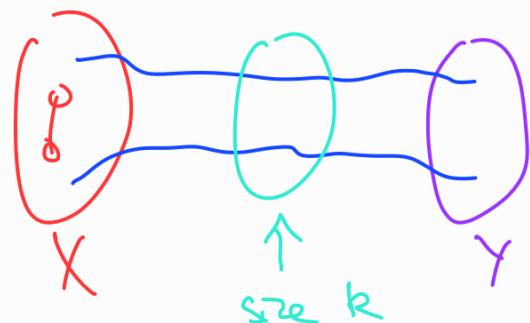
is equal to the maximum number of vertex-disjoint (X, Y) -paths.

Proof continued: Recall, induction on $|E(G)|$. $|E(G)| \geq 1$.

k is the minimum size of a set that separates X from Y

We'd proved

Claim: If G has an edge $e=uv$ with $\{u, v\} \subseteq X$ or $\{u, v\} \subseteq Y$, then there are k vertex-disjoint (X, Y) -paths.



We want to show there are k vertex-disjoint (X, Y) -paths.

We'd picked $e=uv$, and considered G/e (where w is the vertex from the contraction).

Suppose Z' is a set that

separates X' from Y' in G/e

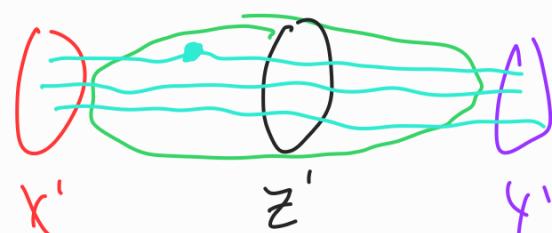
having minimum size.

By induction, there are $|Z'|$ vertex-disjoint (X', Y') -paths.

G



G/e



If $|Z'| = k$, then there are k vertex-disjoint (X, Y) -paths in G .

Suppose $|Z'| < k$.

Case 1: $w \notin Z'$. Then Z' also separates X from Y in G , so

$|Z'| \geq k$, a contradiction.

(Case 2: $w \in Z'$).

Let $Z = (Z' \setminus \{w\}) \cup \{u, v\}$

Then Z separates from Y in G , so

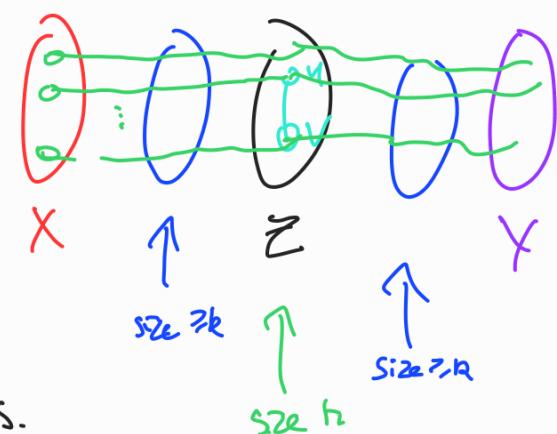
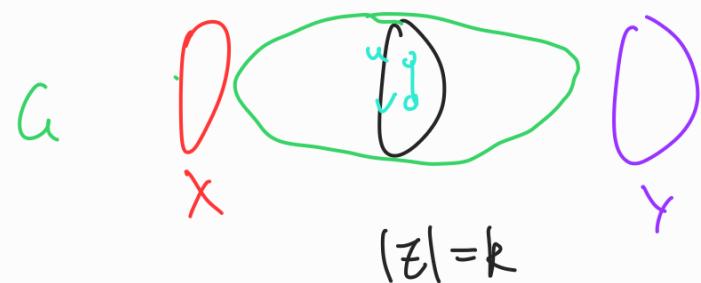
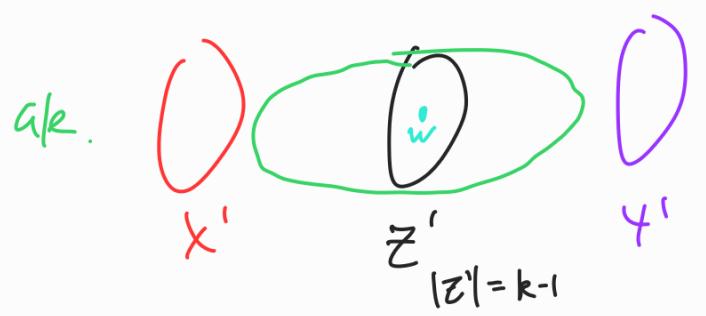
$$|Z| = k.$$

Now, a set that separates X from Z has size at least k (as it would also separate X from Y).

By the earlier claim there are k vertex-disjoint (X, Z) -paths, since $\{u, v\} \subseteq Z$.

Similarly, there are k vertex-disjoint (Z, Y) -paths.

Since $|Z| = k$, by combining these we obtain k vertex-disjoint (X, Y) -paths, as required. \square



Answering the question from yesterday:

there are k vertex-disjoint (X, Y) -paths in G if and only if there is no set of size smaller than k that separates X from Y .

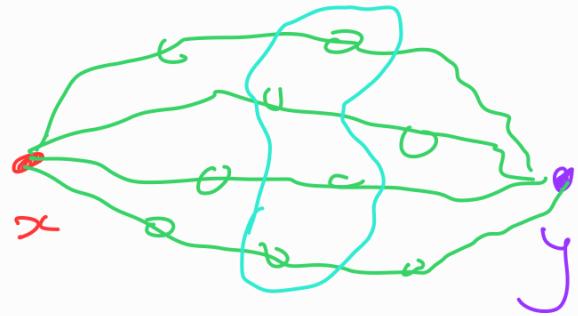
(Corollary 3.29)

Let x and y be distinct vertices in a graph G .

Two (x, y) -paths P and Q

are internally disjoint if

$$V(P) \cap V(Q) = \{x, y\}.$$



When x and y are non-adjacent vertices of G , we

say $S \subseteq V(G) \setminus \{x, y\}$ separates x from y if there
is no path from x to y in $G - S$.

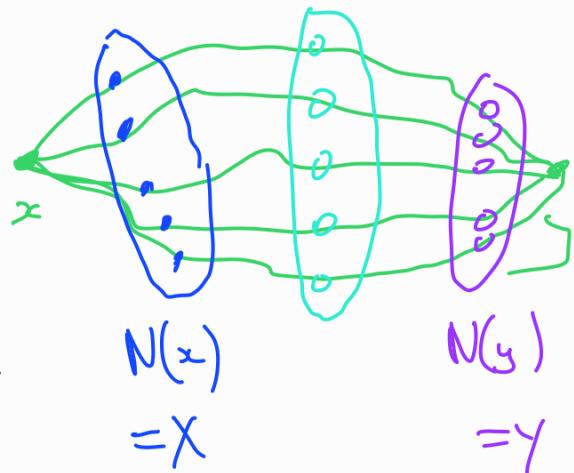
Corollary 3.30: Let G be a graph with non-adjacent
vertices x and y .

The minimum size of a set that separates x and y is
equal to the maximum number of internally disjoint (x, y) -paths.

The neighborhood of a vertex x ,

denoted $N(x)$,

is the set of vertices adjacent to x .



Full details of the proof of

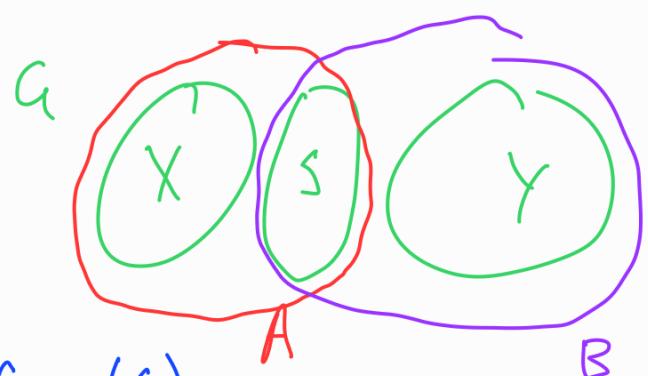
Corollary 3.30 in online notes.

Corollary 3.31: Let G be a graph with more than k vertices, for some positive integer k .

G is k -connected if and only if for every pair of vertices u, v in G , there are k internally disjoint (u, v) -paths.

Separations

Let G be a graph and let A and B be subsets of $V(G)$.

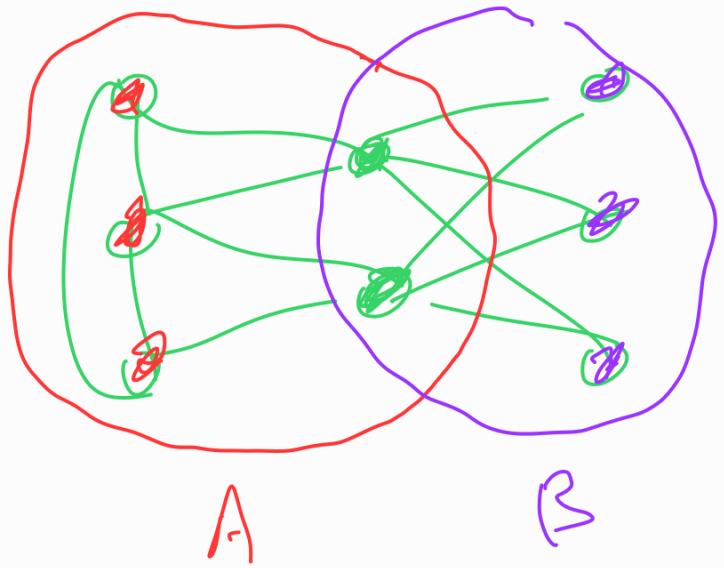


We say $\{A, B\}$ is a separation if

- i) $A \cup B = V(G)$
- ii) there is an edge of G with one end in $A \setminus B$ and the other in $B \setminus A$.

We call $A \cap B$ the boundary of the separation $\{A, B\}$.
and $|A \cap B|$ is the order of the separation

A separation $\{A, B\}$ is proper if both $A \setminus B$ and $B \setminus A$ are non-empty.



$\{A, B\}$ is a proper separation of order 2.